0101
Assembly

ENGR 3410 – Computer Architecture
Fall 2010

What is assembly?

Why are we learning assembly now?

Assembly Language

- Readings: Chapter 2 (2.1-2.6, 2.8, 2.9, 2.13, 2.15), Appendix A.10
- Assembly language
  - Simple, regular instructions – building blocks of C & other languages
  - Typically one-to-one mapping to machine language
- Our goal
  - Understand the basics of assembly language
  - Help figure out what the processor needs to be able to do
- <u>Not</u> our goal to teach complete assembly/machine language programming
  - Floating point
  - Procedure calls
  - Stacks & local variables

MIPS Assembly Language

- The basic instructions have four components:
  - Operator name
  - Destination
  - 1st operand
  - 2nd operand

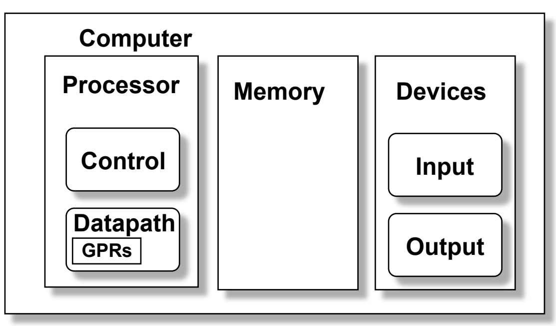
```
add <dst>, <src1>, <src2>     # <dst> = <src1> + <src2>
sub <dst>, <src1>, <src2>     # <dst> = <src1> - <src2>
```

- Simple format: easy to implement in hardware

- More complex: `A = B + C + D − E`

Operands & Storage

- For speed, CPU has 32 general-purpose registers for storing most operands
- For capacity, computer has large memory (64MB+)



- Load/store operation moves information between registers and main memory
- All other operations work on registers

Registers

- 32 registers for operands

Register	Name	Function	Comment
$0	$zero	Always 0	No-op on write
$1	$at	Reserved for assembler	Don't use it!
$2-3	$v0-v1	Function return	
$4-7	$a0-a3	Function call parameters	
$8-15	$t0-t7	Volatile temporaries	Not saved on call
$16-23	$s0-s7	Temporaries (saved across calls)	Saved on call
$24-25	$t8-t9	Volatile temporaries	Not saved on call
$26-27	$k0-k1	Reserved kernel/OS	Don't use them
$28	$gp	Pointer to global data area	
$29	$sp	Stack pointer	
$30	$fp	Frame pointer	
$31	$ra	Function return address	

Basic Operations

(Note: just subset of all instructions)

Mathematic: add, sub, mult, div
Unsigned (changes overflow condition)
Immediate (one input a constant)

```
add $t0, $t1, $t2  # t0 = t1 + t2
addu $t0, $t1, $t2 # t0 = t1 + t2
addi $t0, $t1, 100 # t0 = t1 + 100
```

Logical: and, or, nor, xor
Immediate

```
and $t0, $t1, $t2  # t0 = t1 & t2
andi $t0, $t1, 7   # t0 = t1 & b0111
```

Shift: left & right logical, arithmetic
Immediate

```
sllv $t0, $t1, $t2 # t0 = t1 << t2
sll $t0, $t1, 6    # t0 = t1 << 6
```

Example: Take bits 6-4 of $t0 and make them bits 2-0 of $t1, zeros otherwise:

Memory Organization

- Viewed as a large, single-dimension array, with an address.
- A memory address is an index into the array
- Byte addressing means that the index points to a byte of memory.

0	8 bits of data
1	8 bits of data
2	8 bits of data
3	8 bits of data
4	8 bits of data
5	8 bits of data
6	8 bits of data
...	

Memory Organization (cont.)

- Bytes are nice, but most data items use larger *words*
- For MIPS, a word is 32 bits or 4 bytes.

0	32 bits of data
4	32 bits of data
8	32 bits of data
12	32 bits of data

Our registers hold 32 bits of data

- 2^{32} bytes with byte addresses from 0 to 2^{32}-1
- 2^{30} words with byte addresses 0, 4, 8, ... 2^{32}-4
- Words are aligned
  i.e., what are the least 2 significant bits of a word address?

Endianness

- How do we write numbers?

Must establish a convention for the order of digits to represent numbers

Endianness

- In memory, what is the order of a 32-bit word?

0	32 bits of data
4	32 bits of data
8	32 bits of data
12	32 bits of data

0	8 bits of data
1	8 bits of data
2	8 bits of data
3	8 bits of data

0	
1	
2	
3	

Store the 32-bit word: 0xDEADBEEF

0	32 bits of data
4	32 bits of data
8	32 bits of data
12	32 bits of data

0	8 bits of data
1	8 bits of data
2	8 bits of data
3	8 bits of data

0	
1	
2	
3	

Big and Little Endian

- Big Endian - “Big End” in (first)
  - Motorola 68000
  - Sun SPARC
  - PowerPC G5
  - *Networks*

0	DE
1	AD
2	BE
3	EF

- Little Endian - “Little End” in (first)
  - Intel x86
  - MOS Tech 6502
    - Atari 2600, Apple][, Commodore 64, NES

0	EF
1	BE
2	AD
3	DE

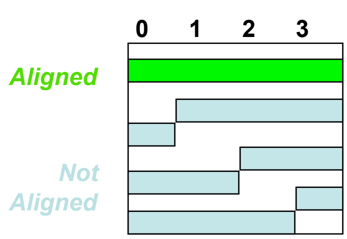
- Bi-Endian - switchable endianness
  - ARM, IBM PowerPC (most)

- Middle-Endian
  - PDP-11

0	
1	
2	
3	

Word Alignment

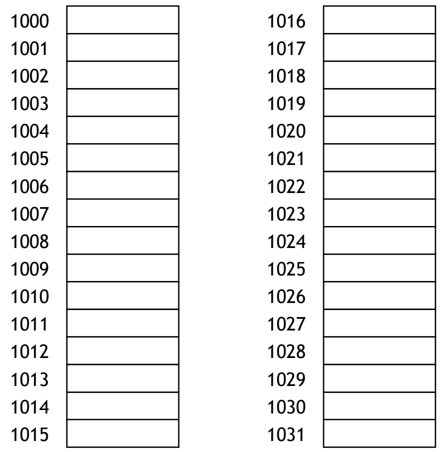
- Require that objects fall on an address that is a multiple of their size



Data Storage

- Characters: 8 bits (byte)
- Integers: 32 bits (word)
- Array: Sequence of locations
- Pointer: Address

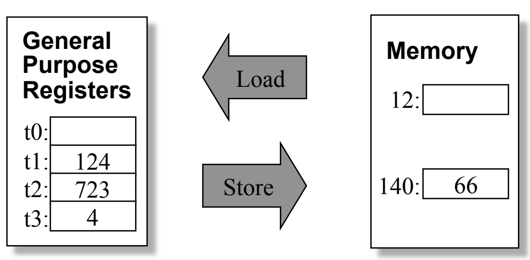
```
char  a = 'G';
int   x = 258;
char *b;
int  *y;
b = new char[4];
y = new int[10];
```



Loads & Stores

- Loads & Stores move data between memory and registers
  - All operations on registers, but too small to hold all data

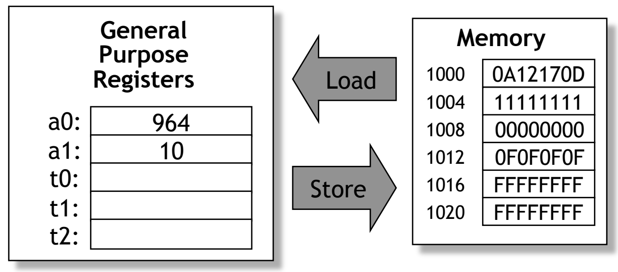
```
lw $t0, 16($t1)      # $t0 = Memory[$t1 + 16]

sw $t2, 8($t3)       # Memory[$t3+8] = $t2
```



- Note: lbu & sb load & store bytes

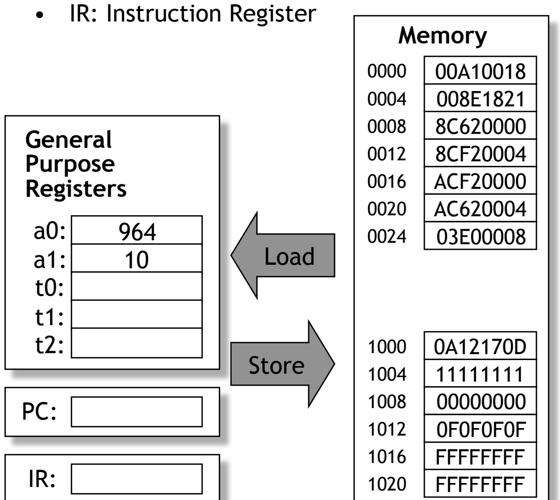
Array Example

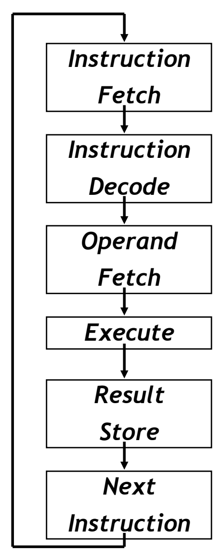
```
/* Swap the kth and (k + 1)th element of an array */
swap(int v[], int k)
{
   int temp = v[k];
   v[k] = v[k + 1];
   v[k + 1] = temp;
}

# Assume v in $a0,
  k in $a1
```



Execution Cycle Example

- PC: Program Counter
- IR: Instruction Register





Control Flow

- Jumps – GOTO different next instruction

```
j 25                 # go to 100: PC = 25 * 4 (instructions are 32-bit)
jr $ra               # go to address in $ra: PC = value of $ra
```

- Branches – GOTO different next instruction if condition is true

  2 register: beq (==), bne (!=)

```
     beq $t0, $t1, FOO# if $t0 == $t1 GOTO FOO: PC = FOO
```

  1 register: bgez (>=0), bgtz (>0), blez (<=0), bltz (<0)

```
     bgez $t0, FOO    # if $t0 >= 0 GOTO FOO: PC = FOO
```

```
if (a == b)
   a = a + 3;
else
   b = b + 7;
c = a + b;
```

```
# $a0 = a, $a1 = b, $a2 = c
     bne   $a0, $a1, ELSEIF     # branch if a!=b
     addi  $a0, $a0, 3;         # a = a + 3
     j DONE;                    # avoid else
ELSEIF:
     addi  $a1, $a1, 7;         # b = b + 7
DONE:
     add   $a2, $a0, $a1;       # c = a + b
```

Loop Example

- Compute the sum of the values 1…N-1

```
int sum = 0;
for (int i = 0; i != N; i++) {
        sum += i;
}
```

```
                         # $t0 = N, $t1 = sum, $t2 = i
```

Comparison Operators

- For logic, want to set a register TRUE (1) / FALSE(0) based on condition

```
slt $t0, $t1, $t2      # if ($t1 < $t2) $t0 = 1 else $t0 = 0;
```

```
if (a >= b)
    c = a + b;              # $t0 = a, $t1 = b, $t2 = c
a = a + c;
```

String toUpper

- Convert a string to all upper case

```
char *index = string;
while (*index != 0) { /* C strings end in 0 */
    if (*index >= 'a' && *index <= 'z')
        *index = *index + ('A' - 'a');
    index++;
}
```

```
# $t0 = index, $t2 = 'a', $t3 = 'z', $t4 = 'A' - 'a', Memory[100] = string
```

Machine Language vs. Assembly Language

- Assembly Language
  - mnemonics for easy reading
  - labels instead of fixed addresses
  - easier for programmers
  - almost 1-to-1 with machine language
- Machine language
  - Completely numeric representation
  - format CPU actually uses

```
SWAP:
    sll   $2, $5, 2                                          000000 00000 00101 00010 00010 000000
    add   $2, $4, $2      // Compute address of v[k]         000000 00100 00010 00010 00000 100000
    lw    $15, 0($2)      // get v[k]                        100011 00010 01111 00000 00000 000000
    lw    $16, 4($2)      // get v[k + 1]                    100011 00010 10000 00000 00000 000100
    sw    $16, 0($2)      // save new value to v[k]          101011 00010 10000 00000 00000 000000
    sw    $15, 4($2)      // save new value to v[k + 1]      101011 00010 01111 00000 00000 000100
    jr    $31             // return from subroutine          000000 11111 00000 00000 00000 001000
```

Labels

- Labels specify the address of the corresponding instruction
  - Programmer doesn't have to count line numbers
  - Insertion of instructions doesn't require changing entire code

```
# $t0 = N, $t1 = sum, $t2 = i
   add   $t1, $zero, $zero   # sum = 0
   add   $t2, $zero, $zero   # i = 0
TOP:
   bne   $t0, $t2, END       # i != N
   add   $t1, $t1, $t2       # sum += i
   addi  $t2, $t2, 1         # i++
   j     TOP                 # next iteration
END:
```

- Notes:
  - Jumps are pseudo-absolute:
    - PC = { PC[31:26], 26-bit unsigned-Address, “00” }
  - Branches are PC-relative:
    - PC = PC + 4 + 4*(16-bit signed Address)

Instruction Types

- Can group instructions by # of operands
  - 3-register
  - 2-register
  - 1-register
  - 0-register

```
add   $t0, $t1, $t2       # t0 = t1 + t2
addi  $t0, $t1, 100       # t0 = t1 + 100
and   $t0, $t1, $t2       # t0 = t1 & t2
andi  $t0, $t1, 7         # t0 = t1 & b0111
sllv  $t0, $t1, $t2       # t0 = t1 << t2
sll   $t0, $t1, 6         # t0 = t1 << 6
lw    $t0, 12($t1)        # $t0 = Memory[$t1 + 10]
sw    $t2, 8($t3)         # Memory[$t3 + 10] = $t2
j     25                  # go to 100 - PC = 25 * 4 (instr are 32-bit)
jr    $ra                 # go to address in $ra - PC = value of $ra
beq   $t0, $t1, FOO       # if $t0 == $t1 GOTO FOO - PC = FOO
bgez  $t0, FOO            # if $t0 >= 0 GOTO FOO - PC = FOO
slt   $t0, $t1, $t2       # if ($t1 < $t2) $t0 = 1 else $t0 = 0;
```

Instruction Formats

- All instructions encoded in 32 bits (operation + operands/immediates)
- Register (R-type) instructions

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 09 08 07 06 05 04 03 02 01 00

OP	RS	RT	RD	SHAMT	FUNCT

(OP = 0,16-20)

- Immediate (I-type) instructions

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 09 08 07 06 05 04 03 02 01 00

OP	RS	RT	16 bit Address/Immediate

(OP = any but 0, 2, 3, 16-20)

- Jump (J-type) instructions

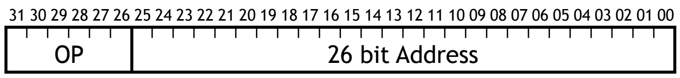
31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 09 08 07 06 05 04 03 02 01 00

OP	26 bit Address

(OP = 2,3)

J-Type

- Used for unconditional jumps

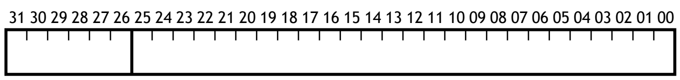


2: j (jump)
3: jal (jump and link)

- Note: top 6 bits of jumped-to address come from current PC

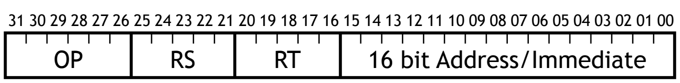
- Example:

```
j    25        # go to 100, PC = 25 * 4 (instr are 32-bit)
```



I-Type

- Used for operations with immediate (constant) operand

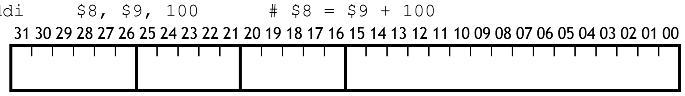


Op1, L/S addr

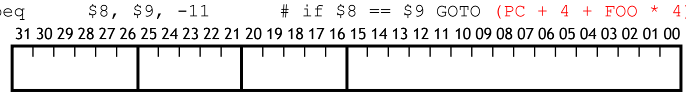
Op2, Dest, L/S targ

04: beq
05: bne
06: blez
07: bgtz
08: addi
09: addiu
10: slti
11: sltiu
12: andi
13: ori
14: xori
32: lb
35: lw
40: sb
43: sw

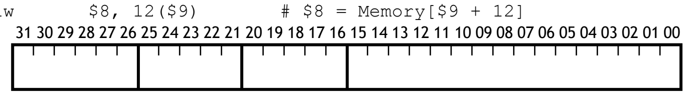
```
addi     $8, $9, 100       # $8 = $9 + 100
```




```
beq      $8, $9, -11       # if $8 == $9 GOTO (PC + 4 + FOO * 4)
```




```
lw       $8, 12($9)        # $8 = Memory[$9 + 12]
```



R-Type

- Used for 3 register ALU operations

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 09 08 07 06 05 04 03 02 01 00

OP	RS	RT	RD	SHAMT	FUNCT
00 (16-20 for FP)	Op1	Op2	Dest	Shift amount (0 for non-shift)	

```
add    $8, $9, $10     # $8 = $9 + $10
```

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 09 08 07 06 05 04 03 02 01 00

```
sll    $8, $9, 6       # $8 = $9 << 6
```

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 09 08 07 06 05 04 03 02 01 00

```
sllv   $8, $9, $10     # $8 = $9 << $10
```

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 09 08 07 06 05 04 03 02 01 00

FUNCT:
00: sll
02: srl
03: sra
04: sllv
06: srlv
07: srav
08: jr
24: mult
26: div
32: add
33: addu
34: sub
35: subu
36: and
37: or
38: xor
39: nor
42: slt

Conversion example

- Compute the sum of the values 0…N-1

```
004   add   $9, $0, $0
008   add   $10, $0, $0
  TOP:
012   bne   $8, $10, END
016   add   $9, $9, $10
020   addi  $10, $10, 1
024   j     TOP
  END:
028
```

Assembly & Machine Language

- Assembly

- Machine Language